Amendments to the Claims

This listing of claims will replace all prior versions, and listings, of claims in the application:

1. (currently amended) A system for organizing and accessing a database, the system comprising:

a secondary <u>B+tree</u> index for <u>indexing</u> a primary B+tree, wherein the secondary <u>B+tree</u> index comprises a plurality of rows each comprising an index key value, and a guess-database address value that represents <u>is</u> a guess as to an <u>what</u> address block of the primary B+tree where a row may be found, where data stored in the database is retrieved using the secondary <u>B+tree</u> index for the primary B+tree.

- (currently amended) The system according to claim 1, wherein the guess-database address values are 4 bytes of address blocks in the primary B+tree.
- 3. (cancelled)
- 4. (currently amended) The system according to claim 1, further comprising:

a guess-database address quality statistic for the secondary index, where the guess-database quess-database address quality statistic represents a ratio of how often the guesses as to where rows may be found in an address block of the primary B+tree are accurate.

5. (currently amended) A <u>computer-implemented</u> method for managing a database system, the method comprising:

creating a secondary <u>B+tree</u> index for <u>indexing</u> a primary B+tree, wherein the secondary B+tree index comprises a plurality of rows each comprising an index key value, and a guess-database address value that represents is a guess as to an what address block of the primary B+tree where a row may be found; and

retrieving data stored in the database system using the secondary <u>B+tree</u> index-for a the primary B+tree.

- 6. (currently amended) The method according to claim 18, further comprising: inserting a row of the secondary index, wherein inserting the row comprises inserting a row comprising an index key value, a mapping table rowid value and a guess database address value
- 7. (previously presented) The method according to claim 18, further comprising: deleting a row of the secondary index, wherein deleting the row comprises locating a row comprising an index key value and a mapping table rowid value and deleting the row.
- updating the secondary index, wherein updating the secondary index comprises locating a row of the secondary index comprising an old index key value and a mapping table rowid value.

8. (currently amended) The method according to claim 18, further comprising:

deleting the row and inserting in the row a new index key value, a mapping table rowid value and a guess database address value.

9. (currently amended) The method according to claim 18, wherein retrieving data stored in the database system further comprises:

obtaining a first guess database address value representing a first address block of the primary B+tree structure:

searching the first address block of the primary B+tree for a row that contains a mapping table rowid value that is the same as a mapping table rowid value in the row where the first guess database address value is stored in the secondary index row; and

if the mapping table rowid is found then the correct row in the primary B+tree has been located and the data is retrieved.

10. (currently amended) The method according to claim 9, wherein retrieving data stored in the database system further comprises:

if the mapping table rowid value is not found, then utilizing the mapping table rowid value stored in the row where the first guess database address value is stored in the secondary index to access a mapping table row stored in a mapping table;

utilizing a second guess-database address value stored in the mapping table row to access a second address block of the primary B+tree;

searching the second address block of the primary B+tree for a primary key that matches a primary key stored in the mapping table row; and

if the primary key is found, then the data is retrieved.

11. (previously presented) The method according to claim 10, wherein if the primary key is not found retrieving data further comprises:

traversing the primary B+tree utilizing the primary key value from the mapping table row to identify the database address to complete the query.

12. (previously presented) The method according to claim 11, further comprising: maintaining a guess-database address quality statistic for the secondary index; maintaining a guess-database address quality statistic for the mapping table; utilizing the statistics to assess guess-database address quality; and

carrying out the query based upon guess-database quality in the secondary index and mapping table, where each of the quess-database address quality statistic represents a ratio of how often the guesses as to where rows may be found in an address block of the primary B+tree are accurate

13. (original) The method according to claim 12, further comprising:

estimating guess-database address quality;

estimating the cost of the query based upon the estimated guess-database address quality; and carrying out the query starting with an index structure with the lowest estimated cost.

14. (currently amended) The method according to claim 5, wherein the guess-database address values are 4 bytes of address blocks in the primary B+tree.

15. (currently amended) A computer program product for performing a process of managing a database system, comprising:

a computer readable medium; and

computer program instructions, recorded on the computer readable medium, executable by a processor, for performing the steps of:

creating a secondary <u>B+tree</u> index for <u>indexing</u> a primary B+tree, wherein the secondary <u>B+tree</u> index comprises a plurality of rows each comprising an index key value, and a guess-database address value that represents <u>is</u> a guess as to an <u>what</u> address block of the primary B+tree where a row may be found; and

retrieving data stored in the database system using the secondary <u>B+tree</u> index-for a the primary B+tree. (currently amended) A system for performing a database management process, comprising

a processor operable to execute computer program instructions; and

a memory operable to store computer program instructions executable by the processor, for performing the steps of:

creating a secondary <u>B+tree</u> index for <u>indexing</u> a primary B+tree, wherein the secondary <u>B+tree</u> index comprises a plurality of rows each comprising an index key value, and a guess-database address value that represents <u>is</u> a guess as to an <u>what</u> address block of the primary B+tree where a row may be found; and

retrieving data stored in the database system using the secondary $\underline{B+tree}$ index for a the primary $\underline{B+tree}$.

17. (previously presented) The system according to claim 1, wherein each row in the plurality of rows further comprising a mapping table rowid value that identifies a row within a mapping table.

18. (previously presented) The method according to claim 5, wherein each row in the plurality of rows further comprising a mapping table rowid value that identifies a row within a mapping table.

19. (previously presented) The computer program product according to claim 15, wherein each row in the plurality of rows further comprising a mapping table rowid value that identifies a row within a mapping table.

20. (previously presented) The system according to claim 16, wherein each row in the plurality of rows further comprising a mapping table rowid value that identifies a row within a mapping table.